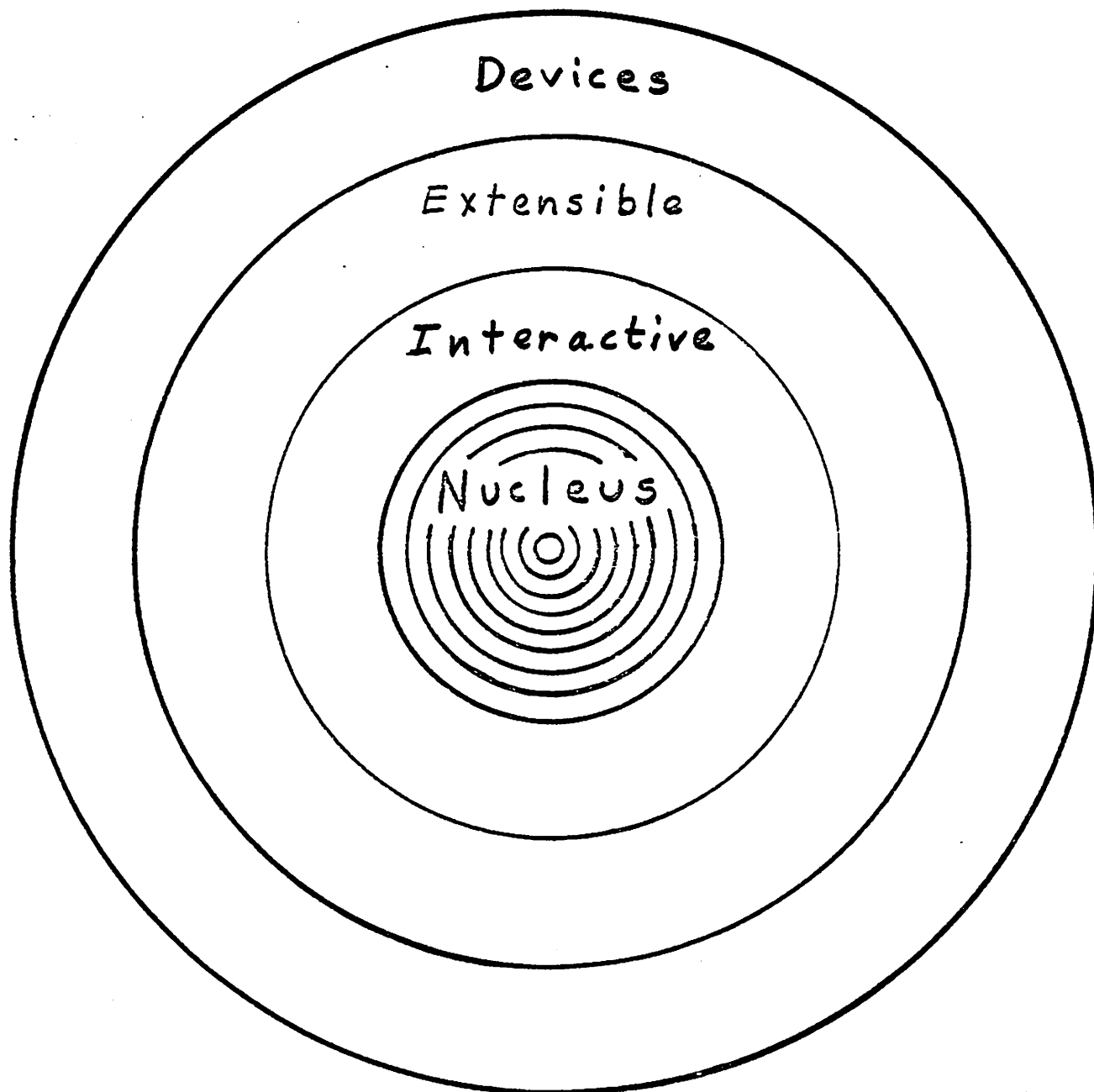


ADDRESS INTERPRETER

Application
Layers



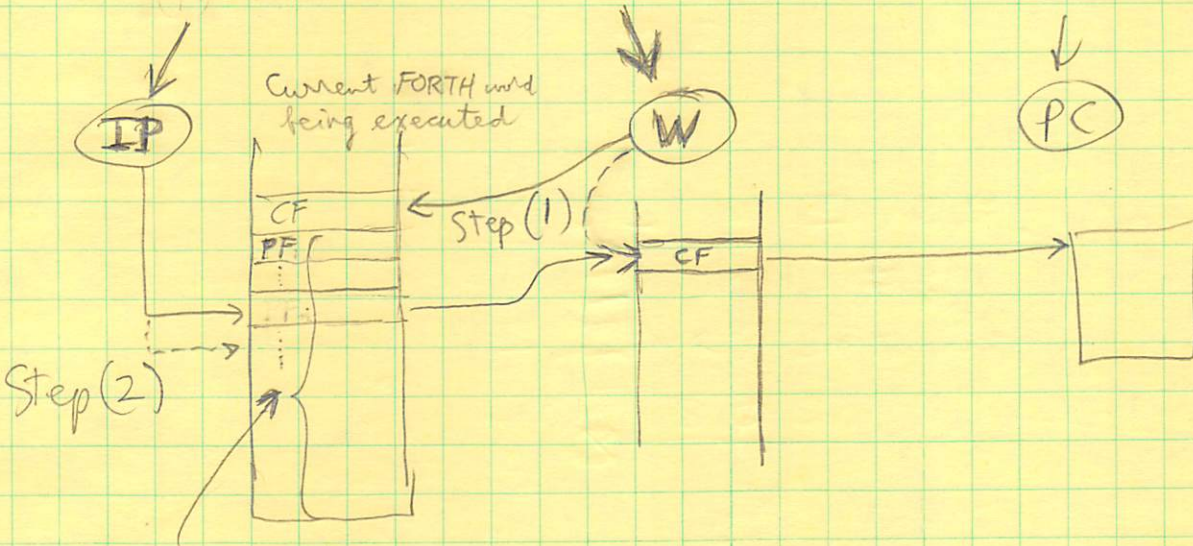
1-12-81

FORTH's address interpreter

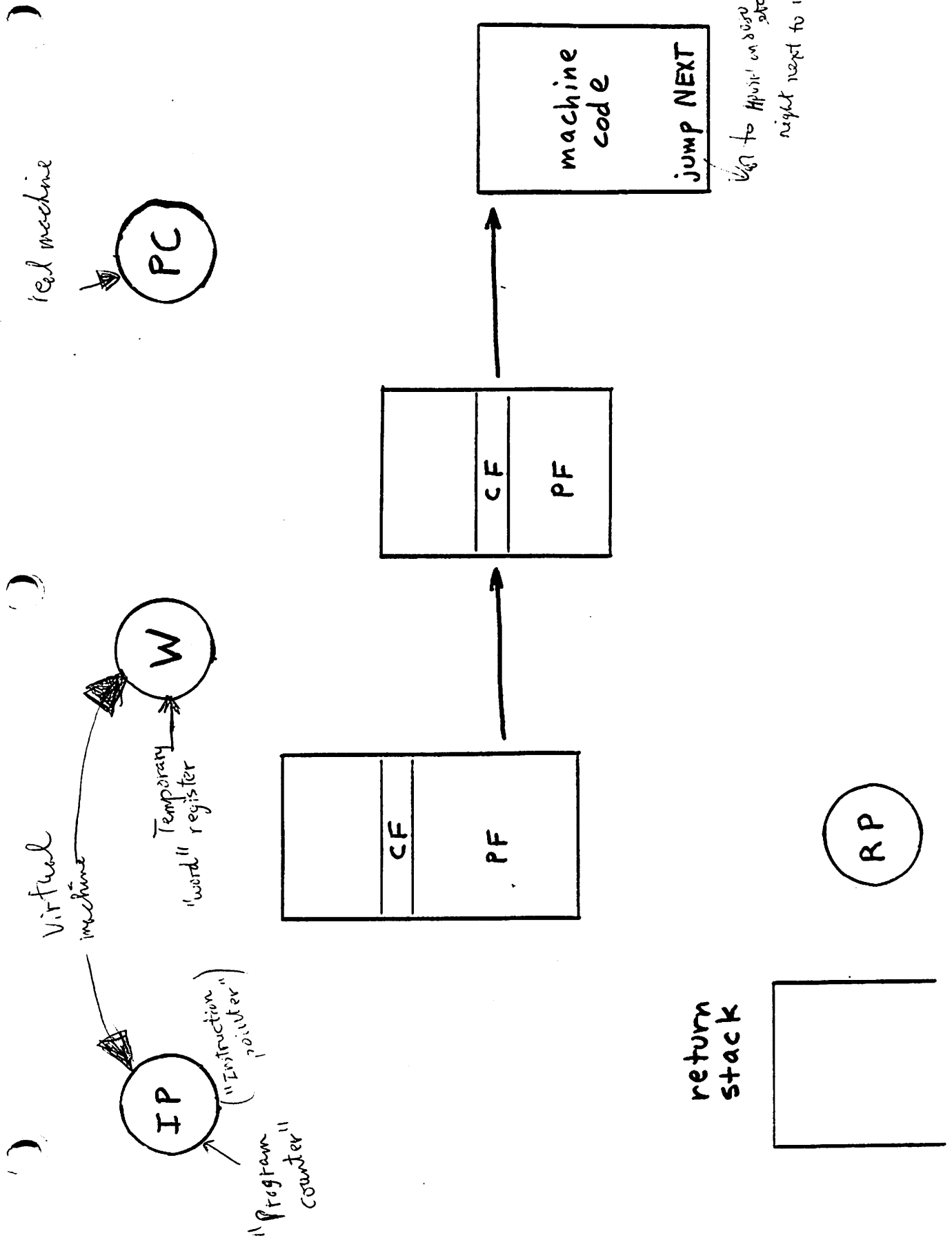
(see Kim Harris course notes p. 108)

FORTH virtual machine

Real machine

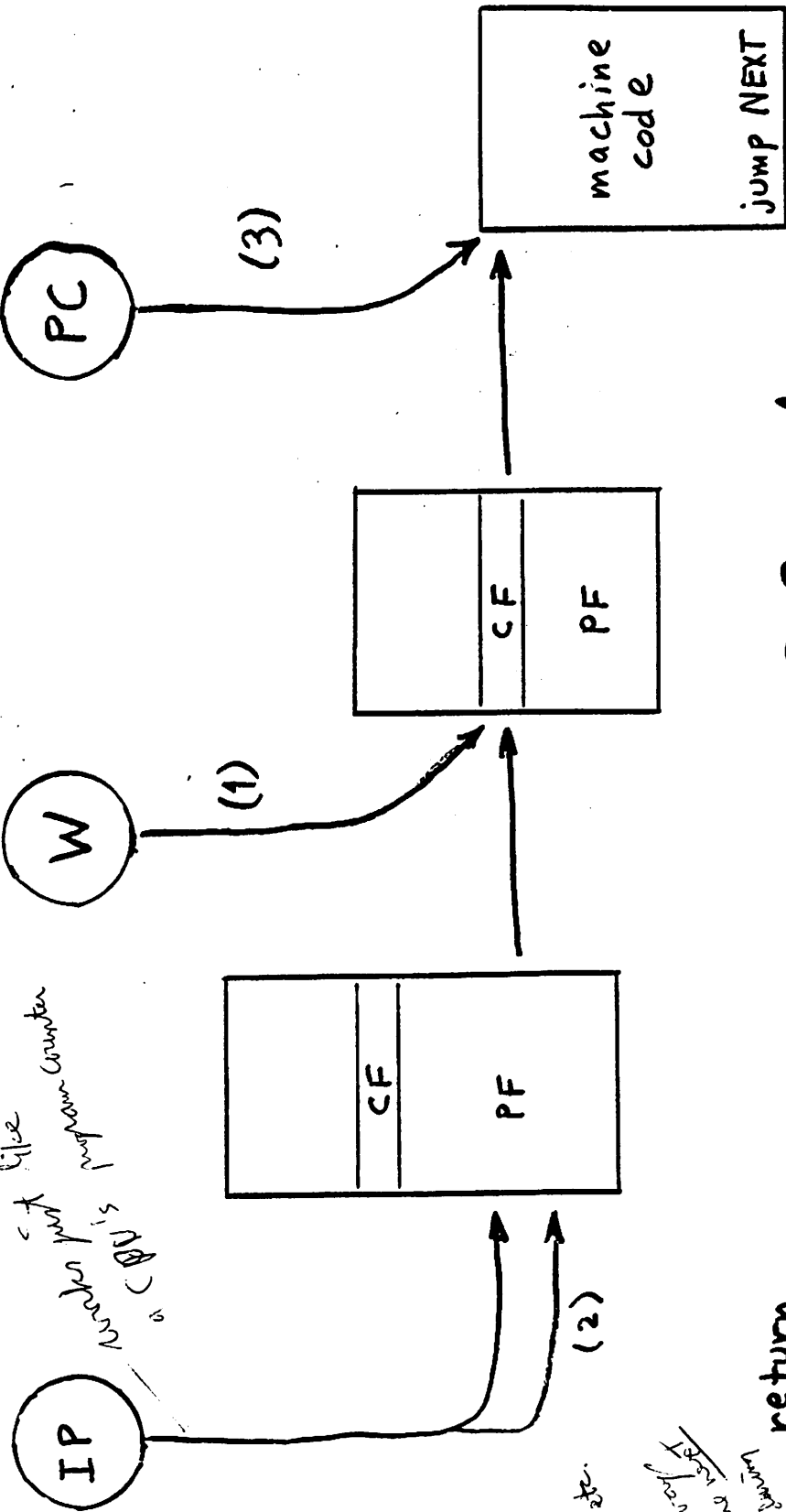


List of Code field addresses
of words used to define this word



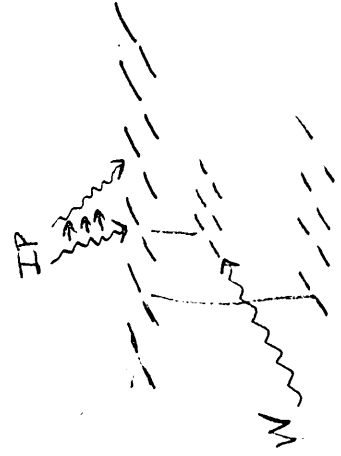
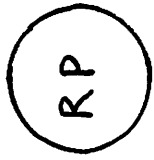
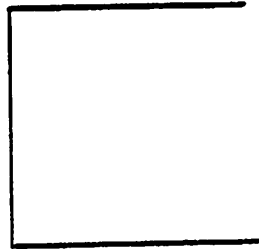
NEXT

FORTH's Address Interpreter



- (1) IP @ W !
- (2) 2 IP +!
- (3) W @ PC !

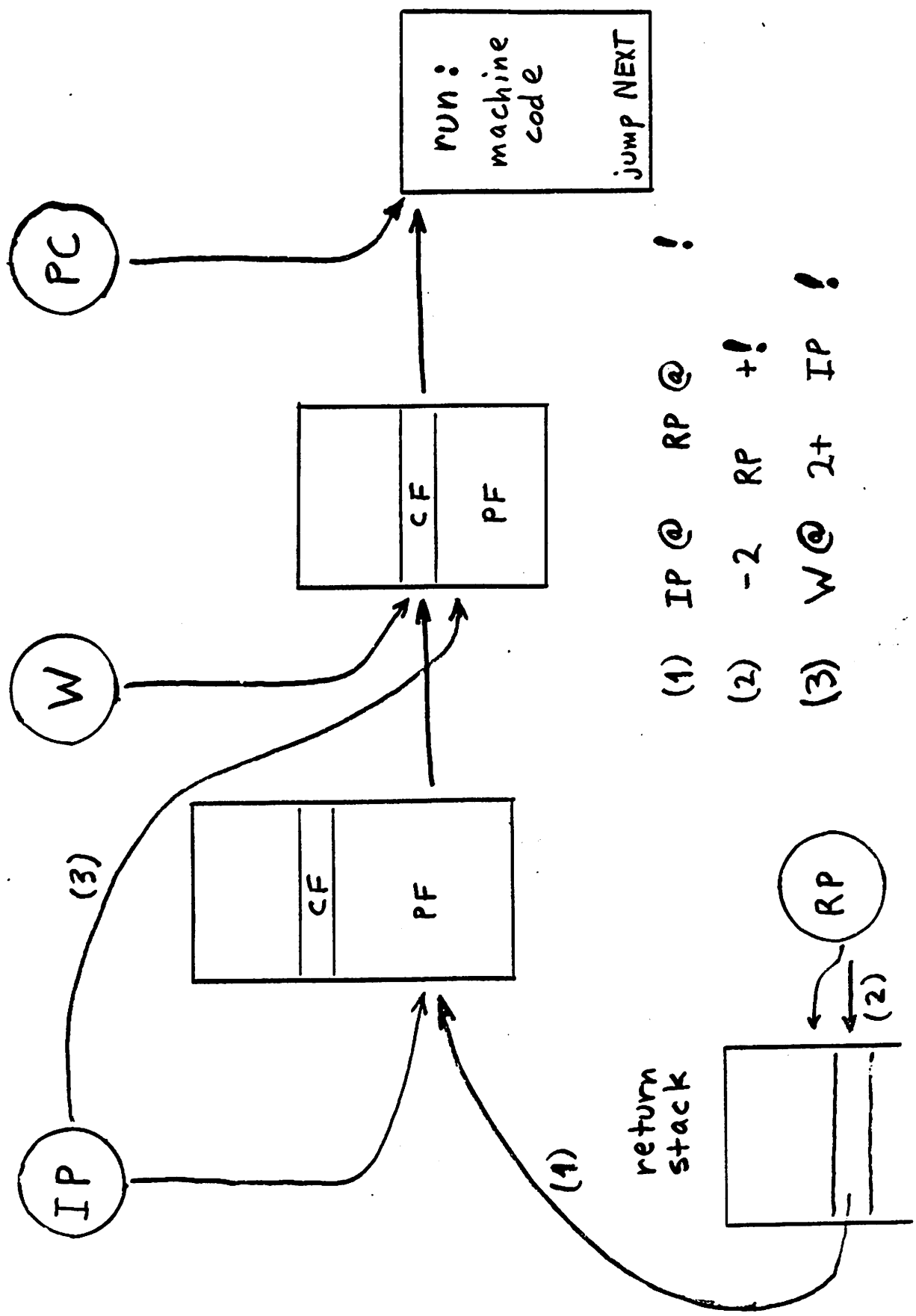
return stack



*like
number pointer
PC's
Program Counter*

*8080 7000
PC
Instruction
return stack
IP
to the next
function
return stack*

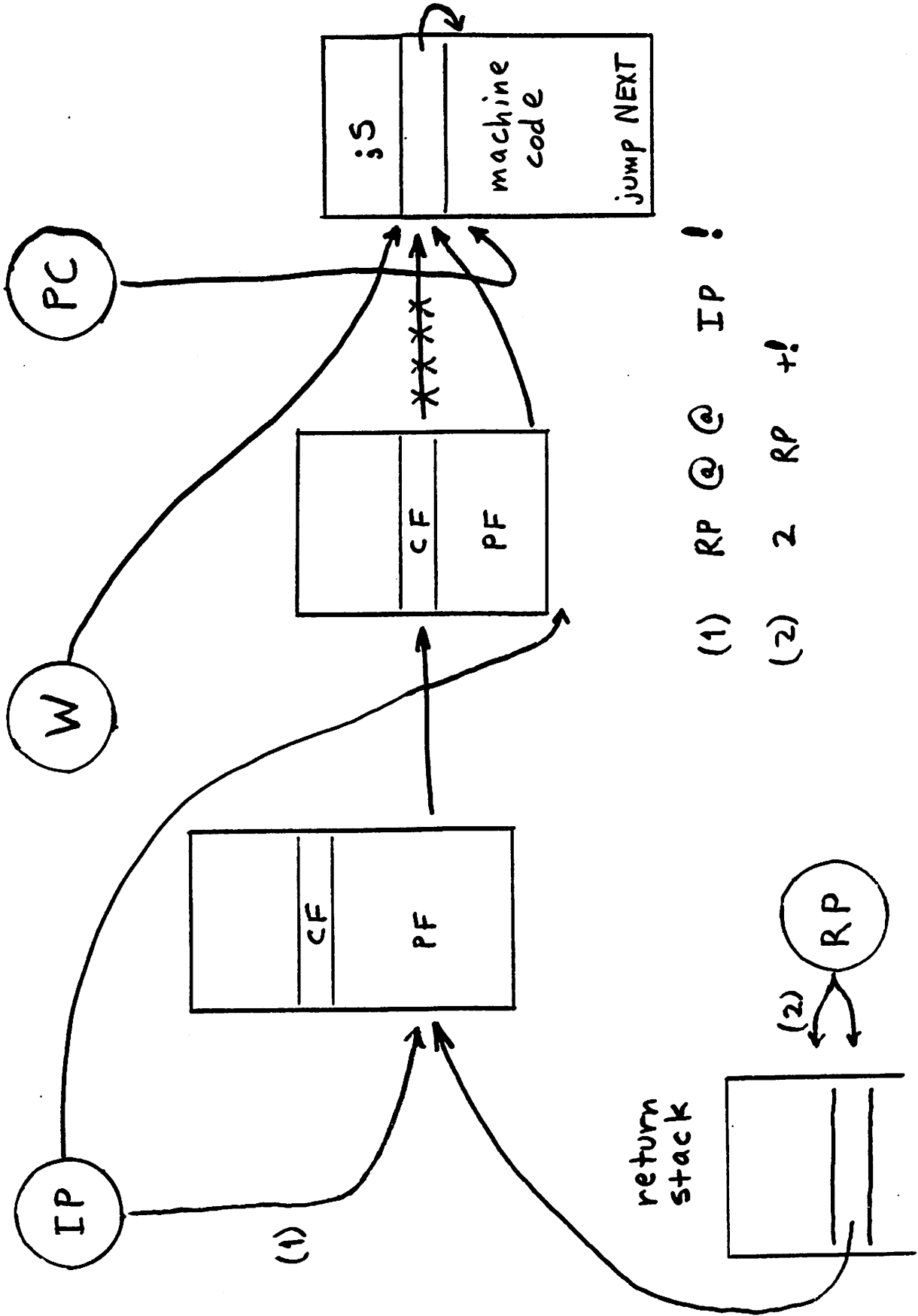
"subroutine" nest



- (1) IP @ RP @ !
- (2) -2 RP +!
- (3) W @ 2+ IP !

"subroutine" un nest

iS



(1) RP @ @ IP !

(2) 2 RP +!

EXECUTION of 8*

IP

W

PC

← dictionary →

INTERPRET

EXECUTE

8*
run:
2*
2*
2*
;S

2*
run:
DUP
+
;S

NEXT
[]

machine code

DUP

JUMP NEXT

+

JUMP NEXT

JUMP NEXT

run:

is

JUMP NEXT

data stack
[]

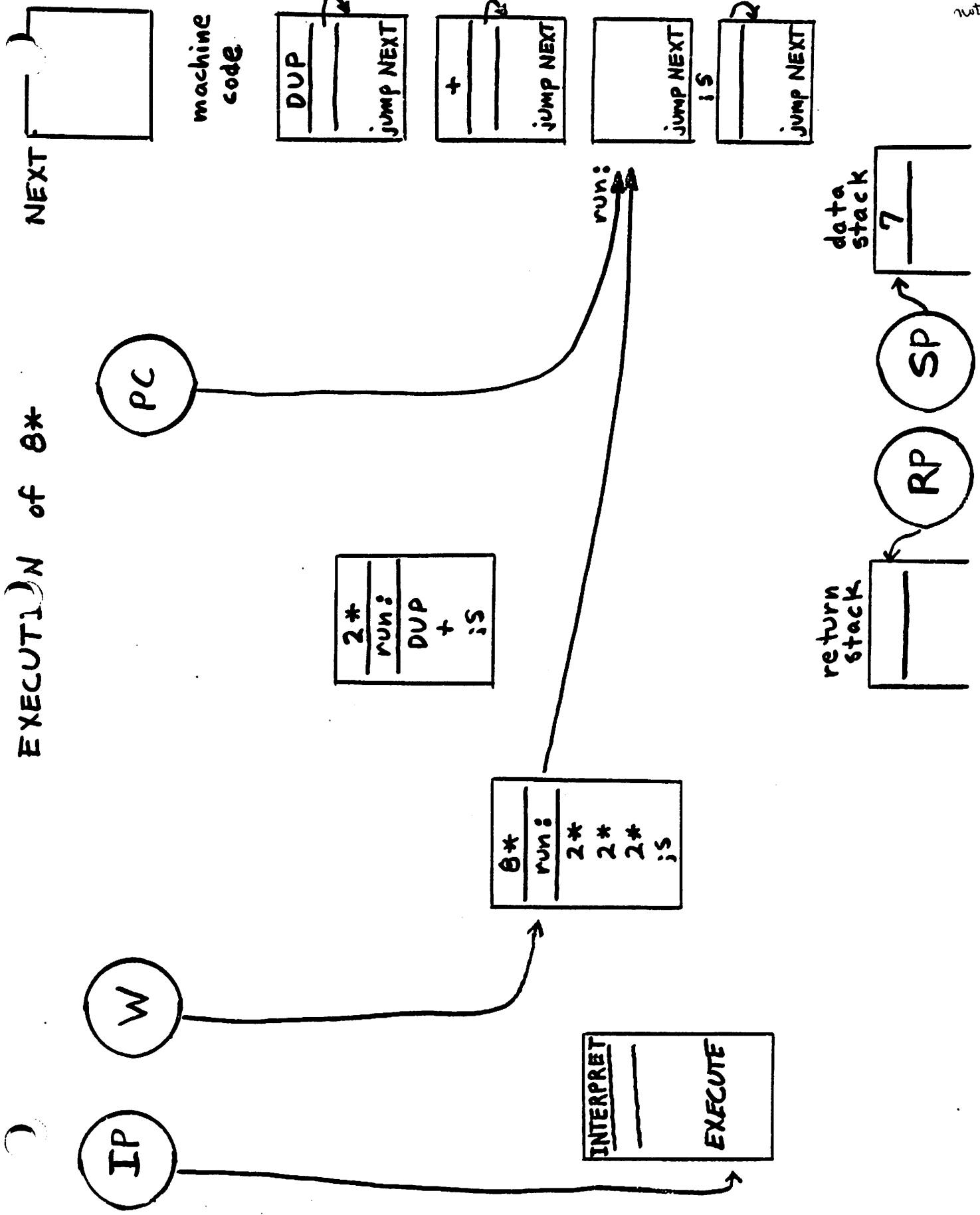
return stack
[]

SP

RP

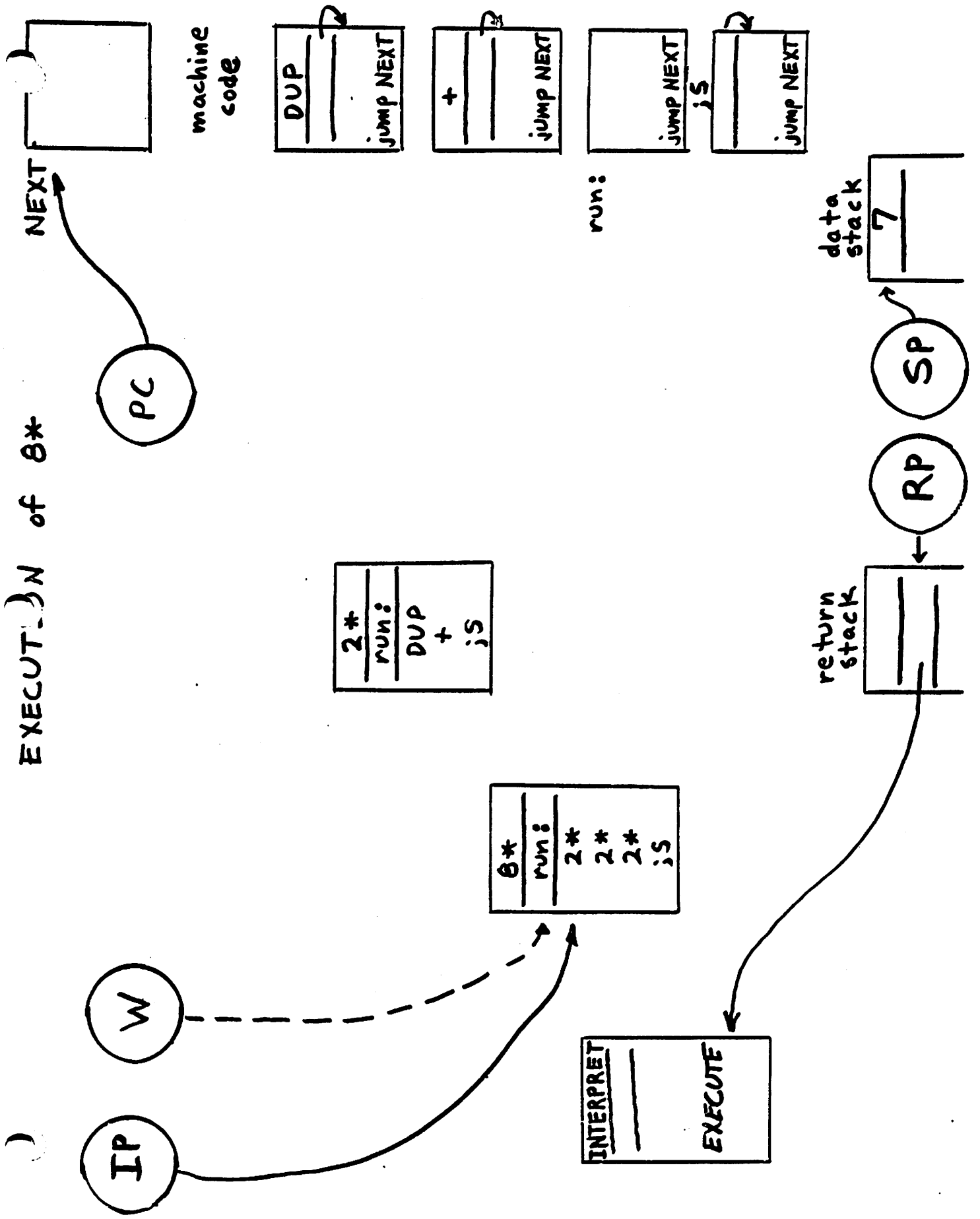
machine code above
use

EXECUTION of 0*

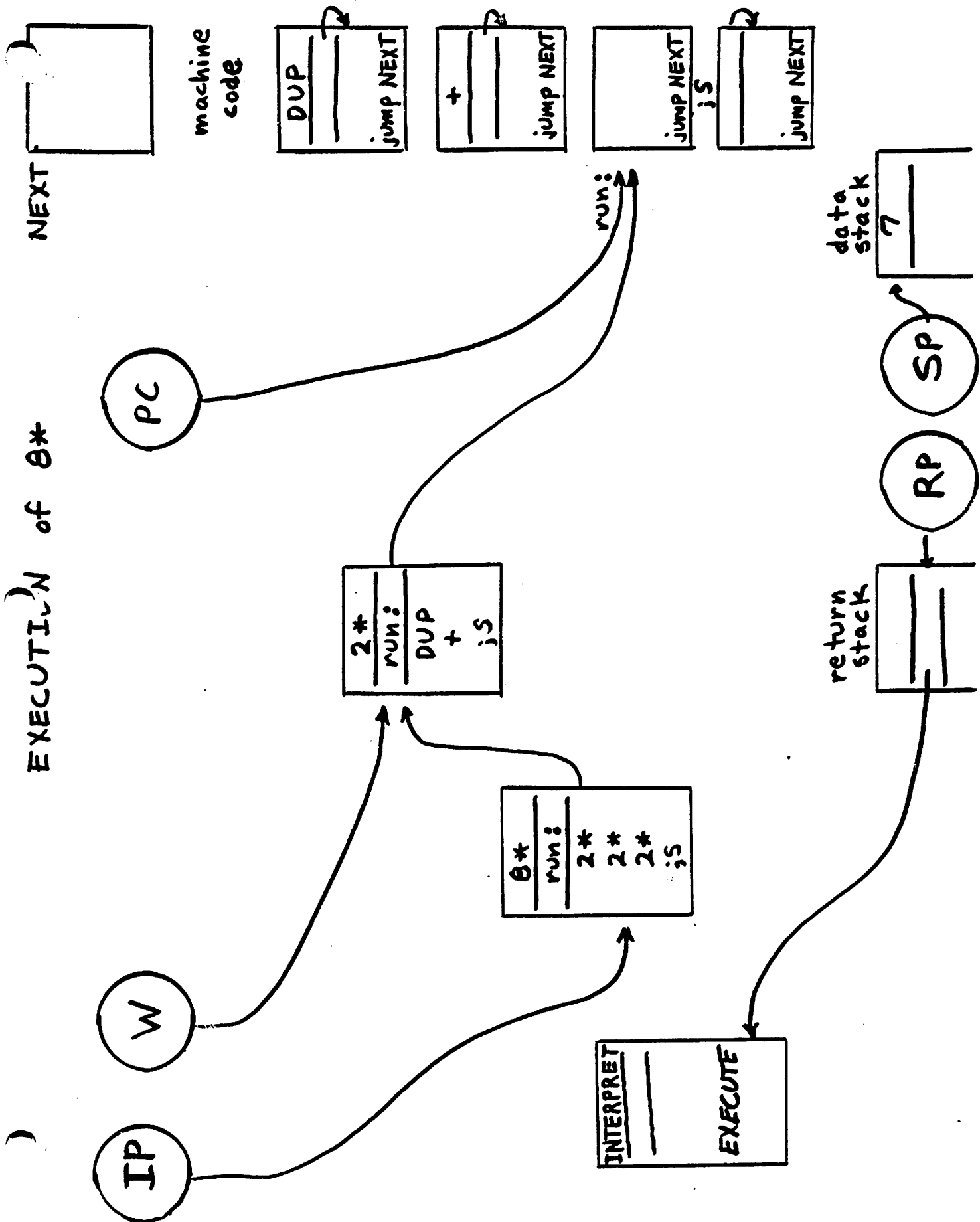


(note: there is wtp. "6")

EXECUTION of 8*

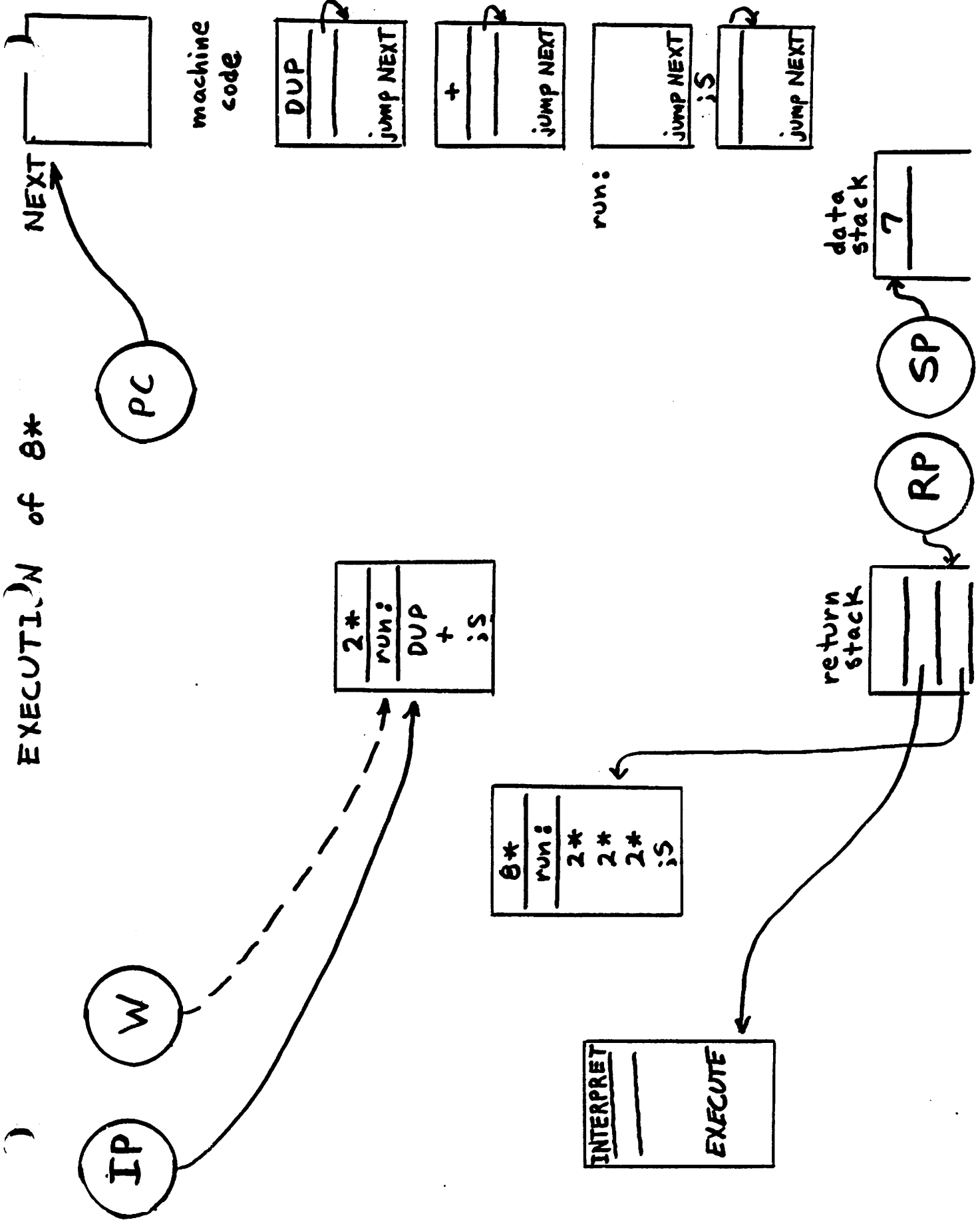


EXECUTION of θ^*



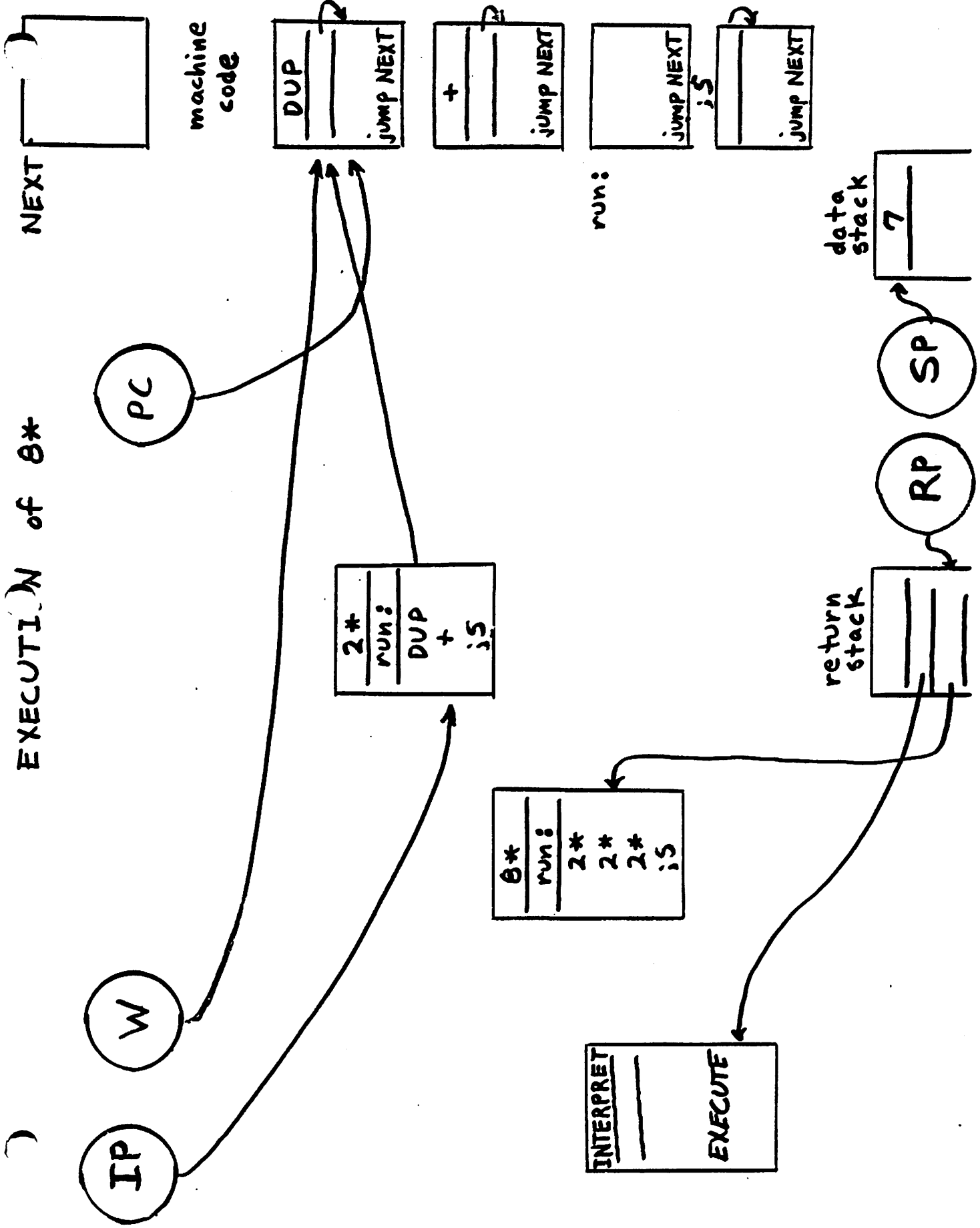
← DICTIONARY →

EXECUTION of 8*



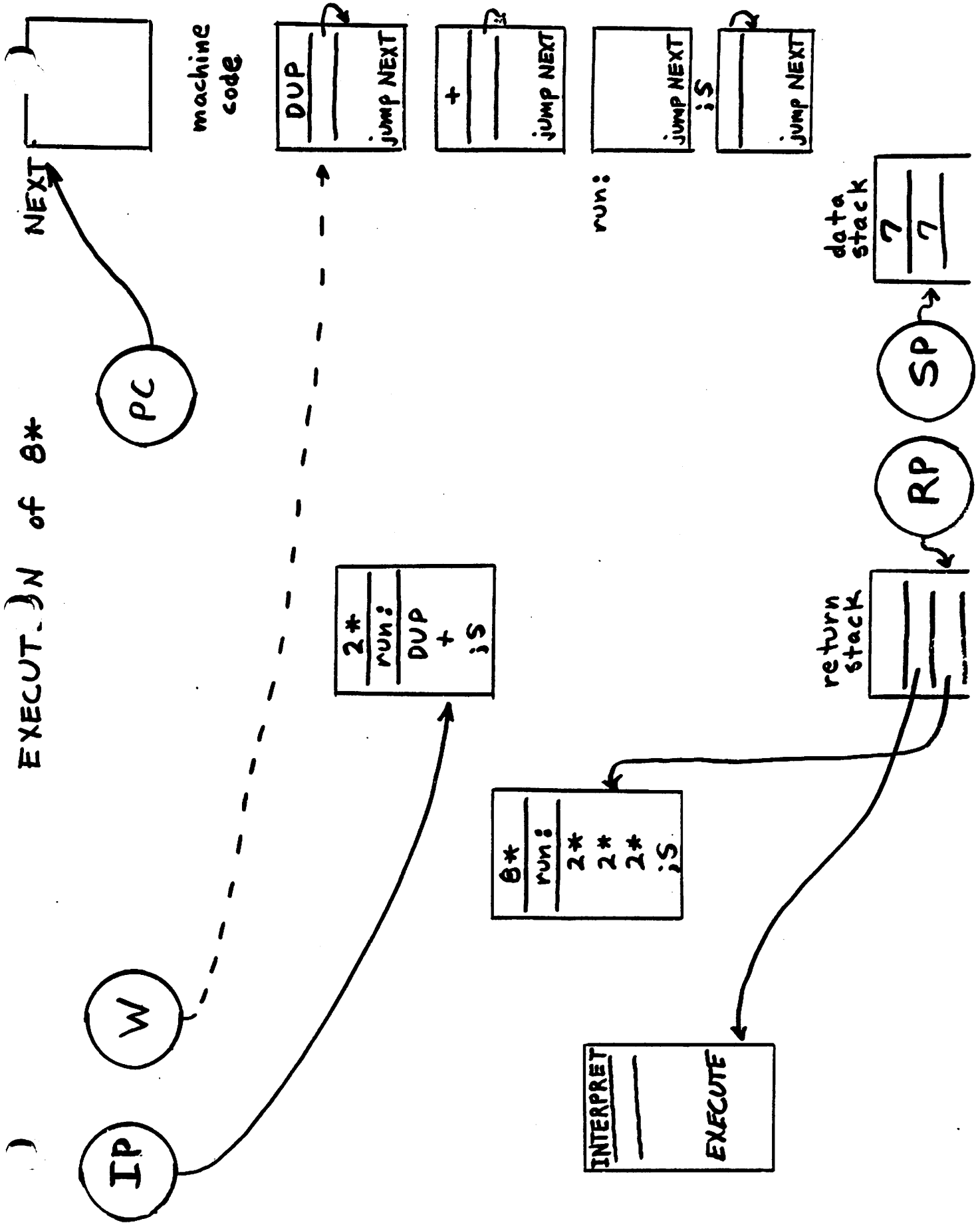
T.S. 4.5

EXECUTION of 0*

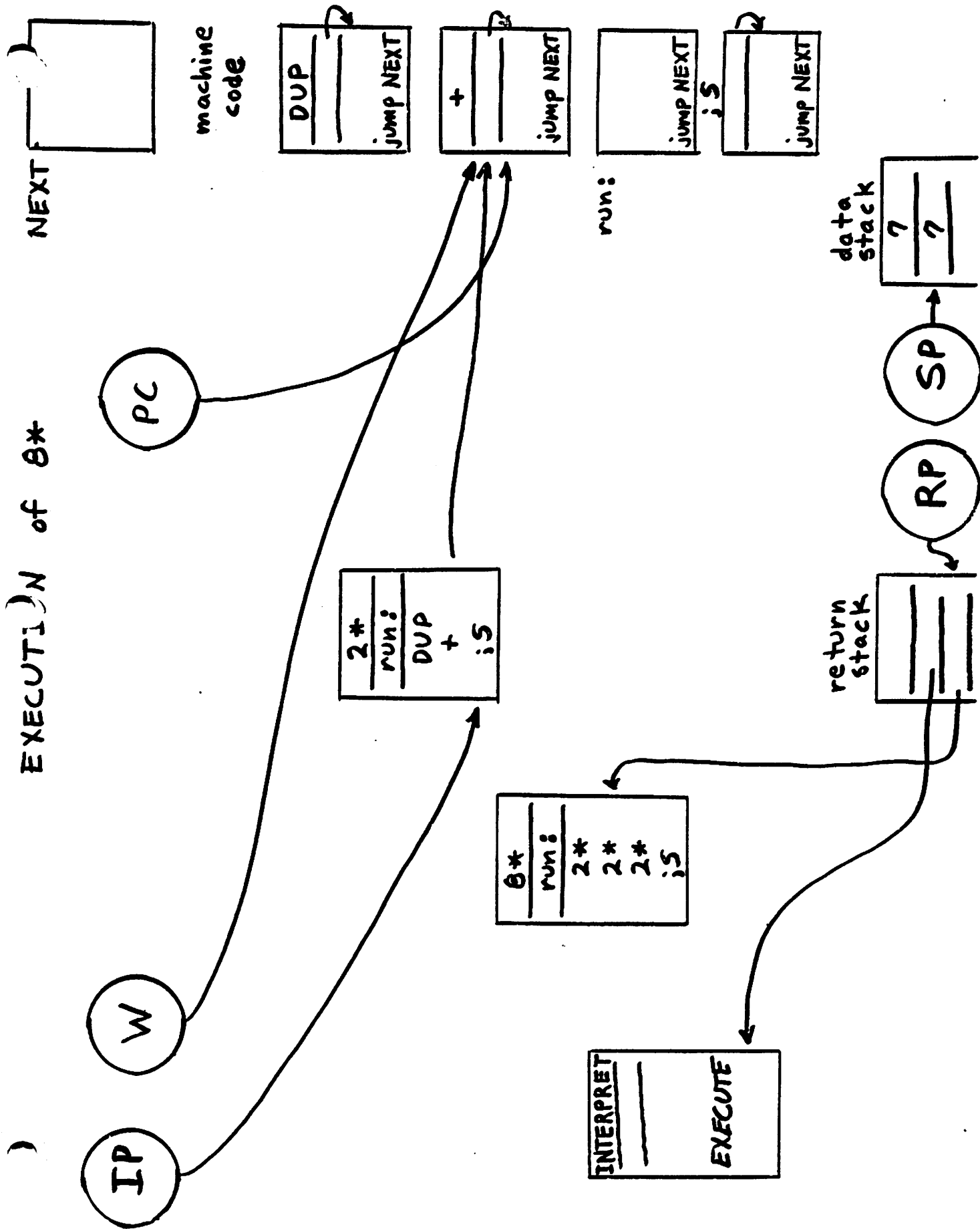


dictionary

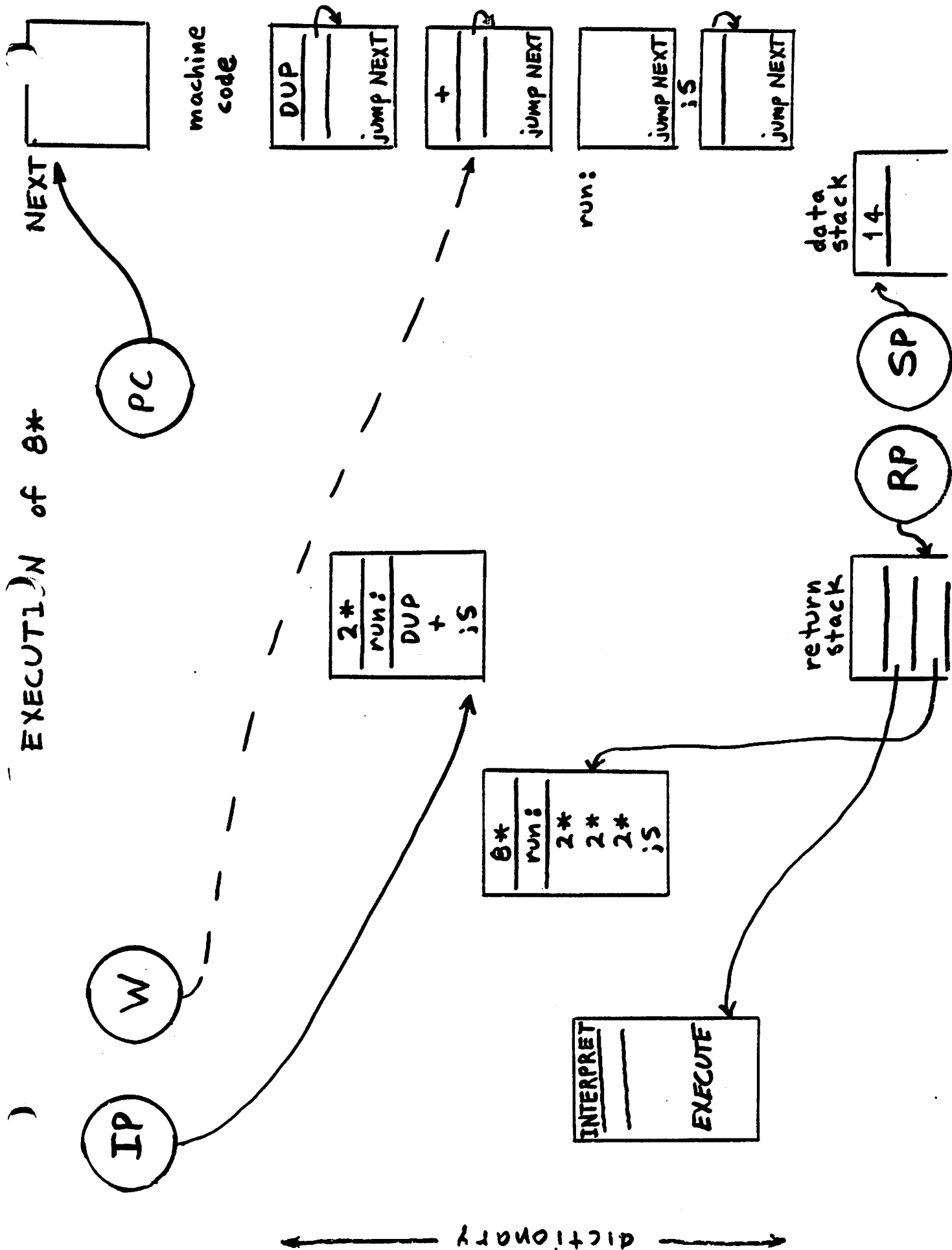
EXECUTION of 8*



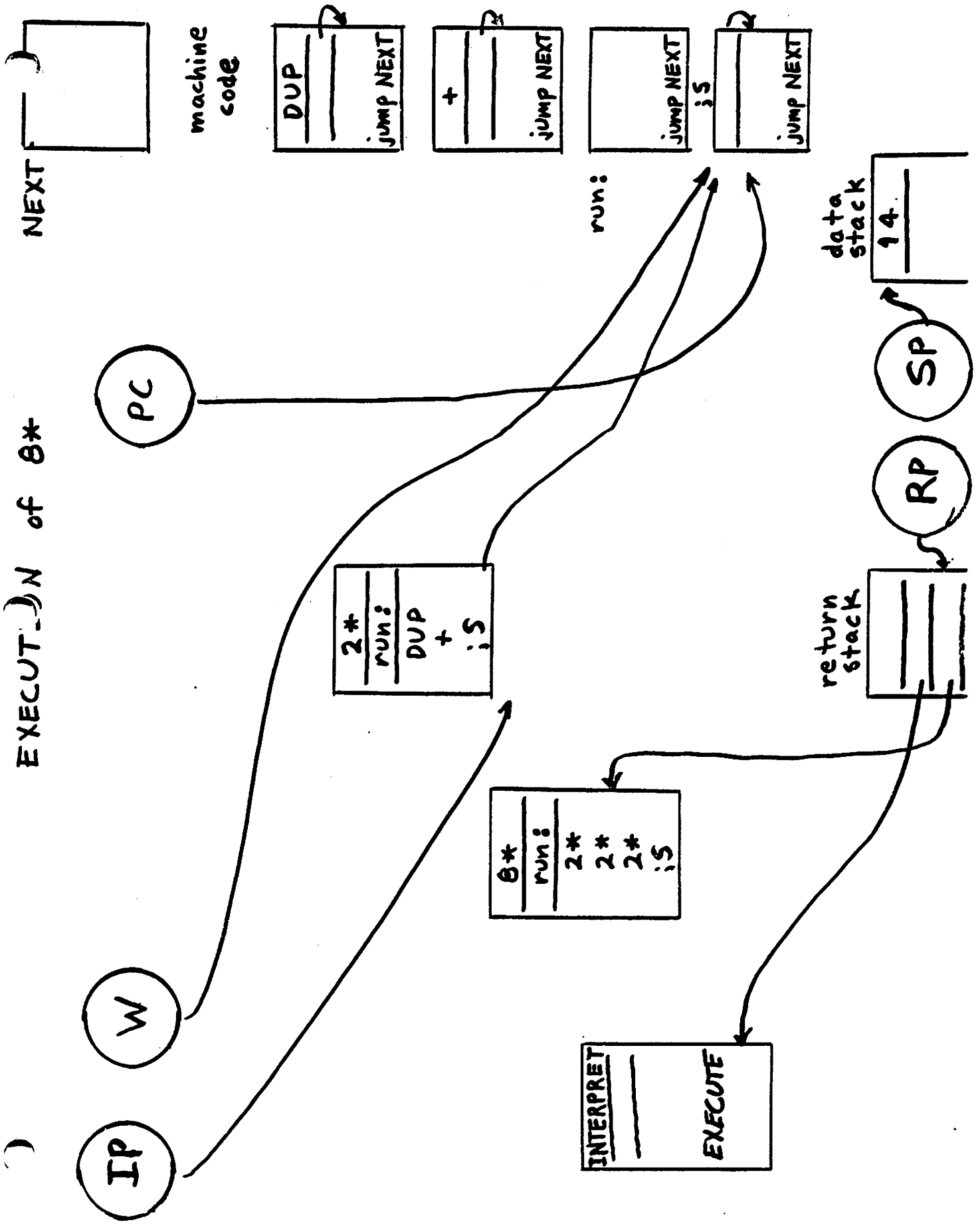
EXECUTION of 0*



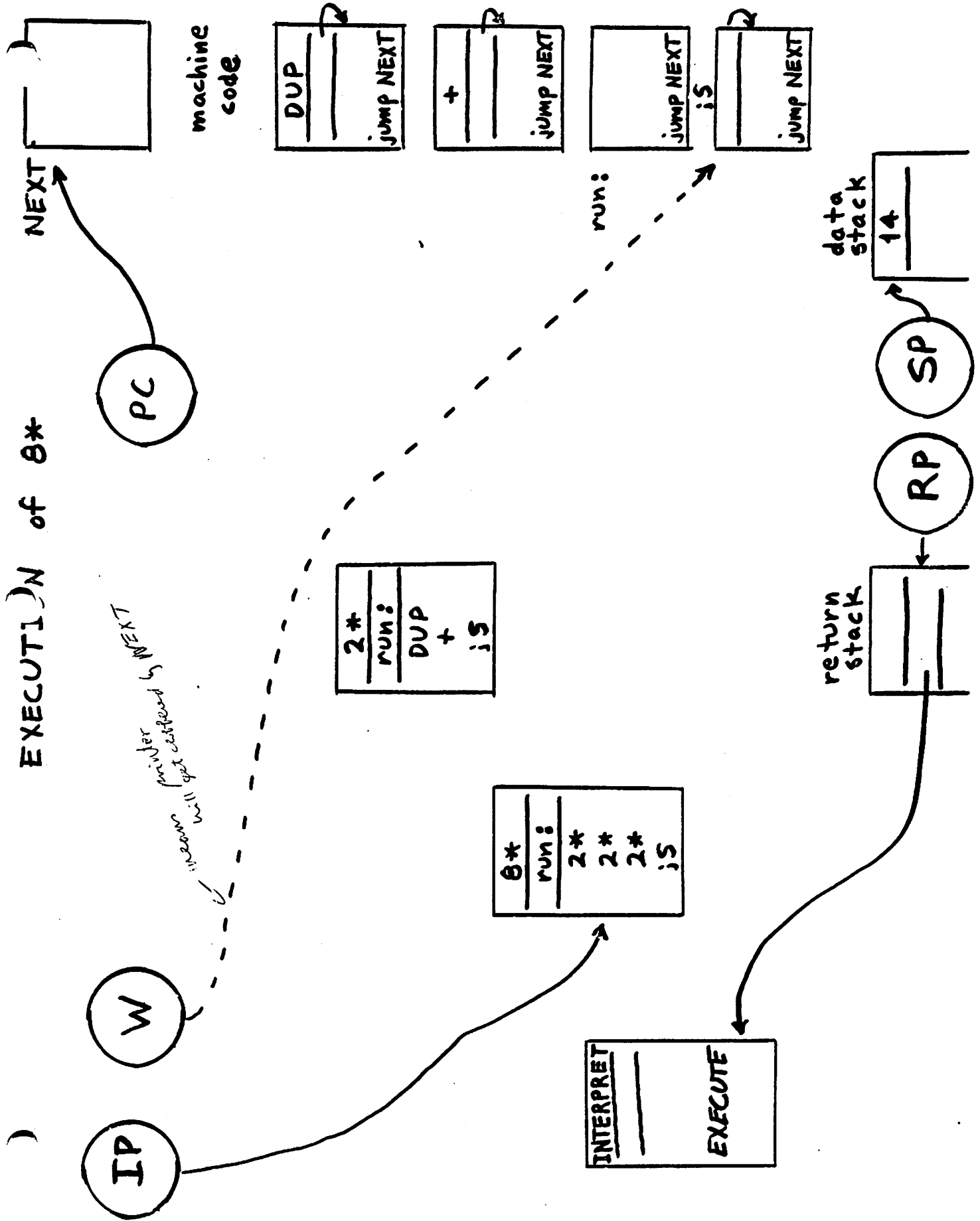
EXECUTION of 8*



EXECUTION of 0*



EXECUTION of 8*



Why is this address interpreter fast inspite of the above overhead? NEXT in 2 instructions on PDP-11

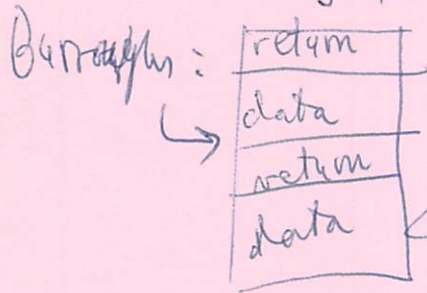
Answer: all we do is fetch addresses

E.g. P-code interpreter of PASCAL has to a detailed set of case statements

microcoded versions of FORTH are running on HP 2100
2109

- may be coming on LXI-11

Why FORTH has 2 stacks



to access this data, requires multiple indirect fetches

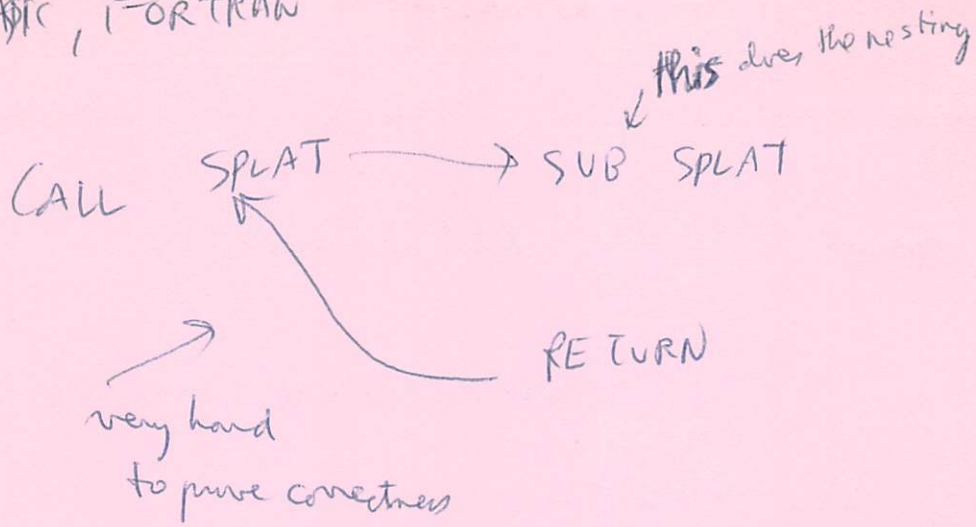
"frame activation & deactivation"
- turns out to be less efficient

one major improvement over ALGOL

Use of W & indirect threaded code allows the nesting operation at the caller

structured nesting & unnesting
(over)

log in BASIC, FORTRAN



in FORTH, we don't have "CALL" → so we can test SPLAT independently

240 B/SCR

* VARIABLE #START

: ELEMENT (index^{subscript} --- addr) 2 * B/BUF/MOD #START
 @ + BLOCK + UPDATE;

(Test virtual array)

: INIT-ARRAY 500 @ DO I I ELEMENT ! LOOP;

: .ARRAY 0 DO CR I. SPACE 2 ELEMENT ? LOOP;

(Make the virtual array into a file)

= AVAILABLE #1 (--- adr) #START @ BLOCK

UPDATE;

(2.9.)

@ AVAILABLE !

(Storing numbers)

3 PUT # AVAILABLE #!

AVAILABLE @ ELEMENT !; @MM

eg. 3 PUT

(Inspecting file entries)

: SEE AVAILABLE @ ^{IT} 1 DO ^{CR I.} I ELEMENT ? LOOP;

Fourth class 6/29/80

2

p. 126

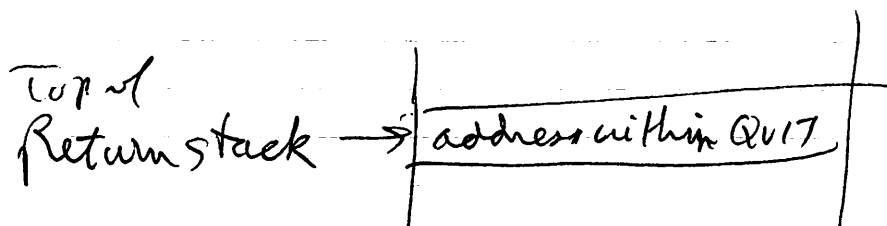
see screen #52

= INTERPRET -FIND

IF STATE @ <

- how to get out (at end of line):

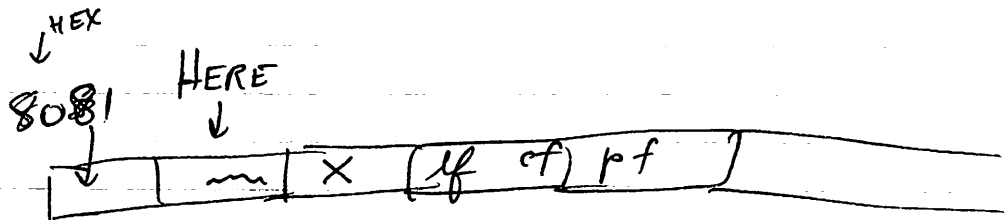
see QUIT



- uses word "null" screen 45

When screen 45 loaded:

8081 is found by interpreter & put on data stack
HERE - executed - - - - - address to



! puts 8081 in addr of HERE

- this puts in a word with name 1 byte long & name is ASCII null

80

- purpose of null: to get us out of whatever we're doing

6/29/80

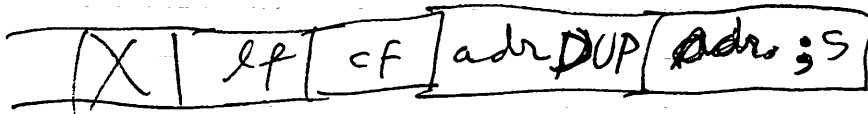
FORTH class

3

Example for p. c9 r1

: X DUP ;

2 ← time frame

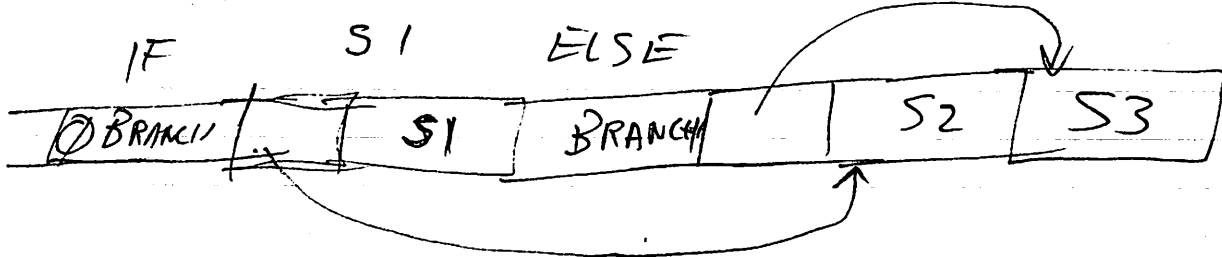


P. C12 r1

NOTE: Data stack not used (by words) ~~by compiler~~ when compiling so available for

P. C14 r1

Desired dictionary result for "ELSE"



CORRECTIONS
revised < & > that work!

< 2DUP XOR OK IF DROP ELSE - THEN OK;

: ~~OK~~ 2DUP XOR OK IF DROP OK 0= ELSE - OK
THEN;

~~SCR #56~~ : S → D DUP OK MINUS ;

SCR #60 : FLUSH #BUF ~~OK~~ IF 0 DO \emptyset
BUFFER DROP LOOP;